

# Mechanism Design for Dynamic Double Auctions

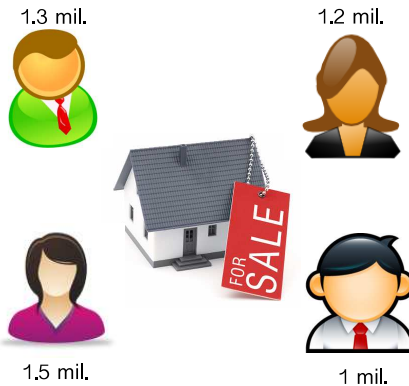
Dengji Zhao<sup>1,2</sup>

<sup>1</sup>ISL, University of Western Sydney, Australia

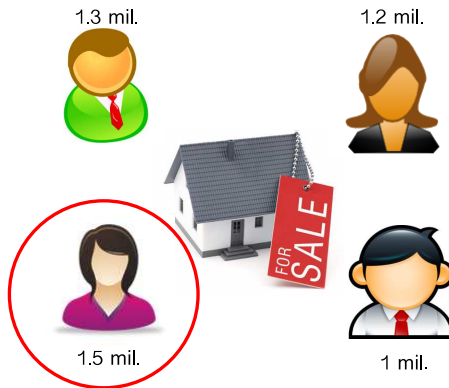
<sup>2</sup>IRIT, University of Toulouse, France

Ph.D. Thesis Defence

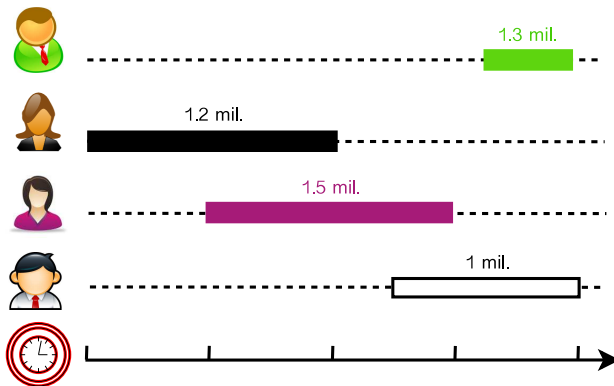
# An Auction



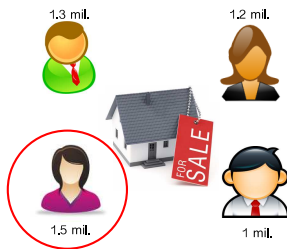
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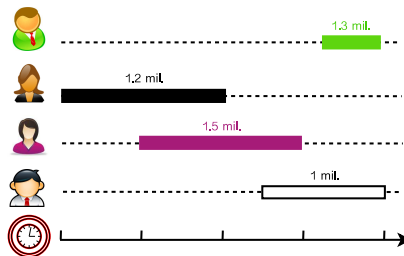
# An Online Auction



# An Online Auction



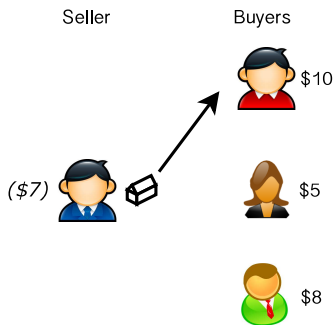
**Static Auction**



**Online Auction**

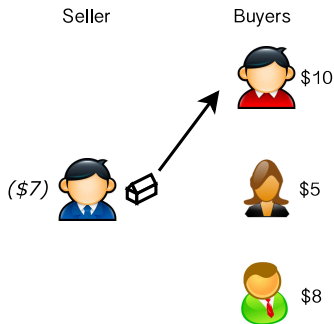
# Double Auction

## One-sided Auction

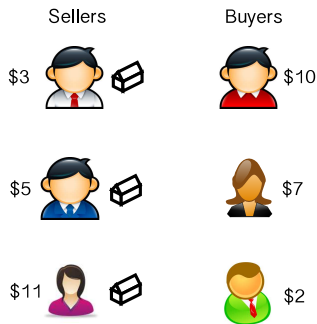


# Double Auction

## One-sided Auction

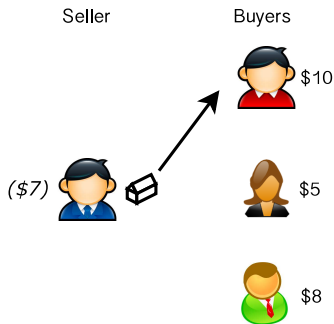


## Double Auction

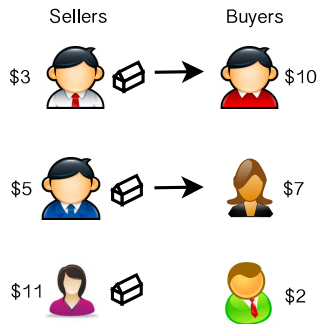


# Double Auction

## One-sided Auction



## Double Auction





# Online Double Auction

## Online Double Auction?

- traders (buyers & sellers) **dynamically** arrive and depart
- traders' valuation (reserve price) **varies** over time

### Real Applications

- Stock Exchange
- Futures Exchange
- Group Buying

# Online Double Auction

## Online Double Auction?

- traders (buyers & sellers) **dynamically** arrive and depart
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- Stock Exchange
- Futures Exchange
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# Online Double Auction

## Online Double Auction?

- traders (buyers & sellers) **dynamically** arrive and depart
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### Real Applications

- Stock Exchange
- Futures Exchange
- Group Buying

# Outline

- 1 Motivation
- 2 The Problem & Solution
- 3 Static Double Auctions
- 4 Approaching Online Double Auctions
- 5 Online Double Auctions
- 6 Conclusion

# Outline

- 1 Motivation
- 2 The Problem & Solution
  - Double Auction Design
  - Design Goals
  - My Solution
- 3 Static Double Auctions
- 4 Approaching Online Double Auctions
- 5 Online Double Auctions
- 6 Conclusion

# The Basic Setting

The roles in a double auction:

- 1 multiple buyers
  - *submit buy orders (called **bids**)*
- 2 multiple sellers
  - *submit sell orders (called **asks**)*
- 3 auctioneer (or market owner)
  - ***match** bids and asks*
  - ***calculate** prices*

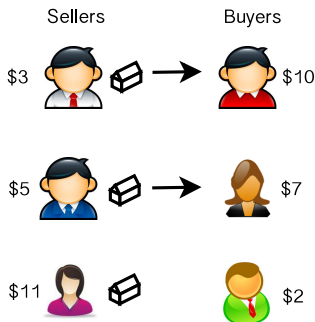
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  - *submit buy orders (called **bids**)*
- 2 multiple sellers
  - *submit sell orders (called **asks**)*
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  - ***calculate** prices*

# What challenge us?

- How to **match** bids and asks?
- How to set exchange **prices**?





# Why are they so challenging?

The decision-making of the auctioneer

- depends on traders' **private information**, called **type**:
  - variable valuation of the commodity/goods
  - dynamical arrival and departure time
  - ...
- has to satisfy certain properties, e.g. truth-telling, efficiency
- faces uncertainty

# Mechanism Design

Mechanism design answers...

How to **design** a mechanism which leads to a **desired outcome**?

A **double auction mechanism** consists of...

- an allocation policy, i.e. a matching between asks and bids
- a payment policy

# What are the Desired Outcomes?

## Question

What properties should a double auction satisfy?

## What are the Desired Outcomes?

## Question

## What properties should a double auction satisfy?

### IC Incentive Compatibility (or Truthfulness)

- telling truth is each trader's **dominant strategy**

### Why IC is important?

- simplifies traders' decision-making
- the base to get other properties, e.g. efficiency

# What are the Desired Outcomes?

## Question

What properties should a double auction satisfy?

### IR Individual Rationality

- no trader receives **negative** utility/profit

*Why IR is important?*

- traders are not forced to participate

# What are the Desired Outcomes?

## Question

What properties should a double auction satisfy?

**Eff.** Efficiency (or Social Welfare Maximisation)

- goods are allocated to the traders who value them **most highly**



# What are the Desired Outcomes?

## Question

What properties should a double auction satisfy?

Liq. Liquidity Maximisation

- (mainly) the number of transactions



# What are the Desired Outcomes?

## Question

What properties should a double auction satisfy?

**Pro.** (Auctioneer's) Profit Maximisation

- **price differences** between matched buyers and sellers

# What are the Desired Outcomes?

## Question

What properties should a double auction satisfy?

**Com.** Computational Complexity

- how much **time** a mechanism takes to compute a solution

# What are the Desired Outcomes?

		Properties						
		IC	IR	Eff.	BB	Liq.	Pro.	Com.
Mechanism (DA)	Allocation (Matching)	X	X	X		X		X
	Payment	X	X		X		X	X

**IC** Incentive Compatibility (or Truthfulness)

**IR** Individual Rationality

**Eff.** Efficiency (or Social Welfare Maximisation)

**BB** Budget Balance (**Weakly** Budget Balance)

**Liq.** Liquidity Maximisation

**Pro.** (Auctioneer's) Profit Maximisation

**Com.** Computational Complexity

# The Problem

**Design Online Double Auctions**, where

- traders dynamically arrive and departure the auction
- traders' valuation is changing over time

## Real Applications

- Stock Exchange
- Futures Exchange
- Group Buying

# An Incremental Approach

- Static
  - single-unit demand/supply with fixed valuation
- Approaching Online (without uncertainty)
  - 1 with temporal constraints, e.g. *futures exchange*
    - temporal constraints limit the allocation/matching
  - 2 under group buying, e.g. *Groupon*
    - valuation varies in terms of the number of goods exchanged
- Online (with uncertainty), e.g. *stock exchange*
  - 1 dynamic arrival & departure + fixed valuation
  - 2 dynamic arrival & departure + dynamic valuation

# The Results

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
	Maximal Matching		X						X	
with temporal constraints	Augmentation-based		X	X	X	X				
under group buying	Existence		X	X	X					
	Impossibility (cond.)		X	X	X					
	Second Price		B	X	X					
	Second Price Plus ...		S	X	X					
	Impossibility		X	X	X		X			
dynamic arrival & departure fixed valuation	Impossibility		X	X						
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based		X		O	O		O		

# Outline

- 1 Motivation
- 2 The Problem & Solution
- 3 Static Double Auctions**
  - The Model
  - The Results
- 4 Approaching Online Double Auctions
- 5 Online Double Auctions
- 6 Conclusion

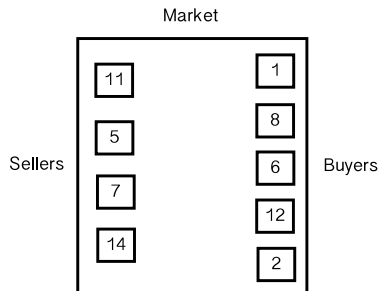
# The Setting

- **type**:  $\theta_i = (v_i)$ 
  - $v_i$  is trader  $i$ 's valuation
- an ask  $\theta_i = (v_i)$  and a bid  $\theta_j = (v_j)$  are **matchable** iff  $v_i \leq v_j$



# The Setting

- traders directly report their types as asks/bids (not necessarily truthfully!)




## The Results

## Equilibrium Matching

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
	Maximal Matching		X						X	
with temporal constraints	Augmentation-based		X	X	X					X
under group buying	Existence		X	X		X				
	Impossibility (cond.)		X	X		X				
	Second Price		B	X		X				
	Second Price Plus ...		S	X		X				
	Impossibility		X	X		X	X			
dynamic arrival & departure fixed valuation	Impossibility		X		X					
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based			X	O		O	O		


# Equilibrium Matching

asks



53
79
85
90
94
98
113
121

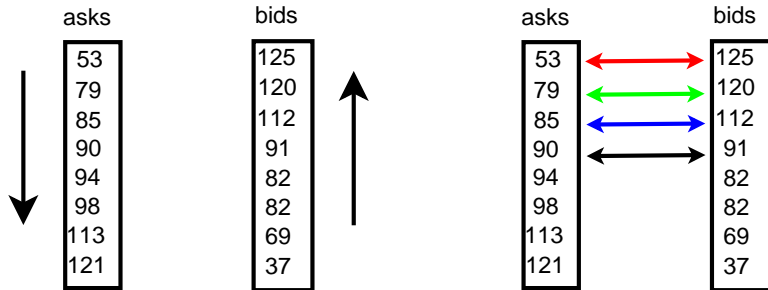
bids



125
120
112
91
82
82
69
37

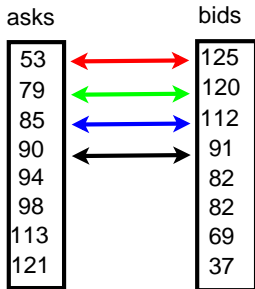
## The Results

## Equilibrium Matching



## The Results

# Equilibrium Matching



## Properties:

- efficient
- (potentially) profit maximization

## The Results

## Equilibrium Matching

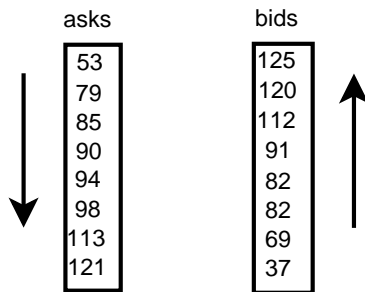
Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
	Maximal Matching		X						X	
with temporal constraints	Augmentation-based		X	X	X					X
under group buying	Existence		X	X		X				
	Impossibility (cond.)		X	X		X				
	Second Price		B	X		X				
	Second Price Plus ...		S	X		X				
	Impossibility		X	X		X	X			
dynamic arrival & departure fixed valuation	Impossibility		X		X					
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based			X	O		O	O		

## The Results

## Maximal Matching

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
	Maximal Matching						X		X	
with temporal constraints	Augmentation-based		X	X	X					X
under group buying	Existence		X	X	X					
	Impossibility (cond.)		X	X	X					
	Second Price		B	X	X					
	Second Price Plus ...		S	X	X					
	Impossibility		X	X	X		X			
dynamic arrival & departure fixed valuation	Impossibility		X	X						
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dynamic arrival & departure dynamic valuation	Behaviour-based		X		O	O		O		

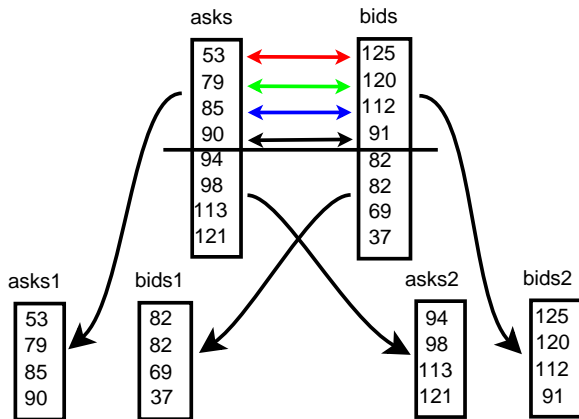
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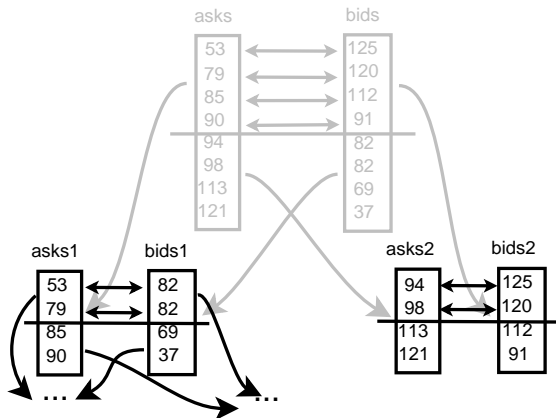


## The Results

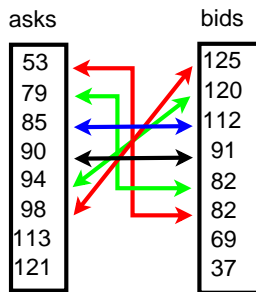
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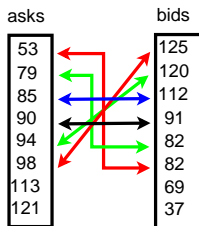


# Maximal Matching



## The Results

## Maximal Matching



## Properties:

- liquidity maximization
- more computationally efficient

## The Results

## Maximal Matching

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
	Maximal Matching						X		X	
with temporal constraints	Augmentation-based		X	X	X					X
under group buying	Existence		X	X	X					
	Impossibility (cond.)		X	X	X					
	Second Price		B	X	X					
	Second Price Plus ...		S	X	X					
	Impossibility		X	X	X		X			
dynamic arrival & departure fixed valuation	Impossibility		X	X						
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based		X		O	O		O		

# Outline

- 1 Motivation
- 2 The Problem & Solution
- 3 Static Double Auctions
- 4 Approaching Online Double Auctions**
  - Double Auction with Temporal Constraints
  - Double Auction under Group Buying
- 5 Online Double Auctions
- 6 Conclusion

# Double Auction with Temporal Constraints

- **type**:  $\theta_i = (v_i, s_i, e_i)$ 
  - $v_i$  is trader  $i$ 's valuation of a single unit of the commodity
  - $s_i$  ( $e_i$ ) is the starting (ending) point of time constraint  $[s_i, e_i]$
- an ask  $\theta_i = (v_i, s_i, e_i)$  and a bid  $\theta_j = (v_j, s_j, e_j)$  are **matchable** iff  $v_i \leq v_j$  and  $[s_i, e_i] \cap [s_j, e_j] \neq \emptyset$

## Application Examples

- Futures Exchange
- Stock Exchange

## Double Auction with Temporal Constraints

## Augmentation-based Mechanism

Environments	Mechanism	Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.
static	Equilibrium Matching		X				O		
	Maximal Matching		X						X
with temporal constraints	Augmentation-based		X	X	X				X
under group buying	<i>Existence</i>		X	X		X			
	<i>Impossibility (cond.)</i>		X	X		X			
	Second Price		B	X		X			
	Second Price Plus ...		S	X		X			
	<i>Impossibility</i>		X	X		X	X		
dynamic arrival & departure fixed valuation	<i>Impossibility</i>		X		X				
	Greedy (cond.)		X	X	X				
	Reduction (cond.)		X	X	X				
dynamic arrival & departure dynamic valuation	Behaviour-based			X	O		O	O	

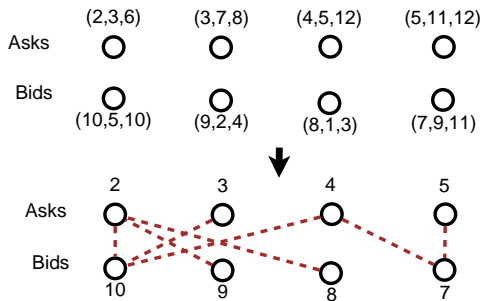


# The Mechanism

- 1 represent asks and bids in a **bipartite graph**
- 2 **maximum-weighted bipartite matching** allocation
- 3 min-max payment

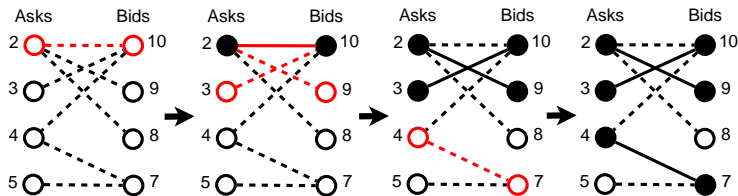
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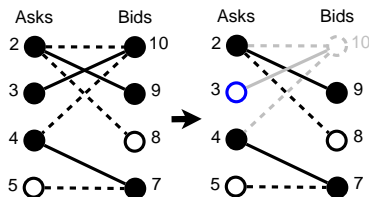
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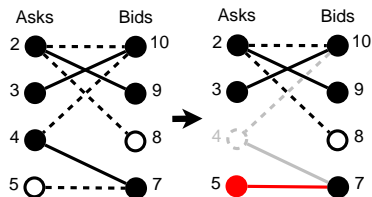


# The Mechanism

- 1 represent asks and bids in a bipartite graph
- 2 maximum-weighted bipartite matching allocation
- 3 min-max payment



the lowest price to win



the highest price to win

# The Properties

- truthful, efficient, individually rational (i.e. VCG mechanism)
- complexity
  - can be implemented  $O(n)$  times faster than the classical VCG mechanism

## Double Auction with Temporal Constraints

## Augmentation-based Mechanism

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
	Maximal Matching		X						X	
with temporal constraints	Augmentation-based		X	X	X	X				
under group buying	Existence		X	X	X					
	Impossibility (cond.)		X	X	X					
	Second Price		B	X	X					
	Second Price Plus ...		S	X	X					
	Impossibility		X	X	X			X		
dynamic arrival & departure fixed valuation	Impossibility		X	X						
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based		X			O	O			O

Double Auction under Group Buying

# Group Buying



**GROUPON**  
Collective Buying Power

  
livingsocial.

# The Setting

Trader  $i$  has valuation function  $v_i : \mathbb{Z} \rightarrow \mathbb{R}$ .

- Seller:

- unlimited supply
- **monotonic**:  $v_i(k) \leq v_i(k+1)$
- **group buying discount**:  $\frac{v_i(k)}{k} \geq \frac{v_i(k+1)}{k+1}$

- Buyer:

- demands  $c_i$  units
- $v_i(k) = v_i(c_i) > 0$  for all  $k \geq c_i$ , otherwise  $v_i(k) = 0$



## Double Auction under Group Buying

## Existence of IC, IR, BB Auctions

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
	Maximal Matching		X						X	
with temporal constraints	Augmentation-based		X	X	X	X				
under group buying	Existence		X	X	X					
	Impossibility (cond.)		X	X	X					
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	Second Price Plus ...		S	X	X					
	Impossibility		X	X	X		X			
dynamic arrival & departure fixed valuation	Impossibility		X	X						
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based		X		O	O		O		

# Existence of IC, IR, BB Auctions

## Existence Example

- **fixed-price auctions**, i.e. price doesn't depend on traders



## Double Auction under Group Buying

## Impossibility I

Environments	Mechanism	Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.
static	Equilibrium Matching		X				O		
	Maximal Matching		X						X
with temporal constraints	Augmentation-based		X	X	X				X
under group buying	<i>Existence</i>		X	X		X			
	<i>Impossibility (cond.)</i>		X	X		X			
	Second Price		B	X		X			
	Second Price Plus ...		S	X		X			
	<i>Impossibility</i>		X	X		X	X		
dynamic arrival & departure fixed valuation	<i>Impossibility</i>		X		X				
	Greedy (cond.)		X	X	X				
	Reduction (cond.)		X	X	X				
dynamic arrival & departure dynamic valuation	Behaviour-based			X	O		O	O	

# Impossibility I

given that

- both the trading size and the payment are neither seller-independent nor buyer-independent

there is no auction that is incentive compatible, individually rational and (weakly) budget-balanced.

Why?

buyers want to form a bigger group while sellers might not!

- buyers with larger group will lower their payments
- a seller's profit might not maximised when selling more

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# Partially Truthful Auctions

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# Partially Truthful Auctions

## Second Price Auction

Given type profile report  $v = (v^B, v^S)$ , assume that  $v_1^B(1) \geq v_2^B(1) \geq \dots \geq v_m^B(1)$ .

- 1 Let  $w(k) = \min \arg \min_i v_i^S(k)$  and  $p(k) = \min_{i \neq w(k)} \frac{v_i^S(k)}{k}$  or  $\infty$  if there is only one seller.
- 2 Let  $k^* = \max\{k \mid v_k^B(1) \geq p(k)\}$ .
- 3 The first  $k^*$  buyers, i.e. buyers of valuation  $v_1^B, v_2^B, \dots, v_{k^*}^B$ , receive one unit of the commodity each and each of them pays  $p(k^*)$ .
- 4 Seller  $w(k^*)$  sells  $k^*$  units of the commodity and receives payment  $p(k^*) \cdot k^*$ .
- 5 The remaining traders lose without payment.





## Double Auction under Group Buying

# Impossibility II

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
	Maximal Matching		X						X	
with temporal constraints	Augmentation-based		X	X	X	X				
under group buying	Existence		X	X	X					
	Impossibility (cond.)		X	X	X					
	Second Price		B	X	X					
	Second Price Plus ...		S	X	X					
	Impossibility		X	X	X		X			
dynamic arrival & departure fixed valuation	Impossibility		X	X						
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based		X		O	O		O		

# Impossibility II

- There is no (weakly) budget balanced, individually rational, truthful auctions that can guarantee trading size.

## Double Auction under Group Buying

## Impossibility II

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
	Maximal Matching						X		X	
with temporal constraints	Augmentation-based		X	X	X					X
under group buying	Existence		X	X	X					
	Impossibility (cond.)		X	X	X					
	Second Price		B	X	X					
	Second Price Plus ...		S	X	X					
	Impossibility		X	X	X		X			
dynamic arrival & departure fixed valuation	Impossibility		X	X						
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based		X		O	O		O		

# Outline

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- 5 Online Double Auctions**
  - with Fixed Valuation
  - with Dynamic Valuation
- 6 Conclusion

## Online Double Auctions

- 1 dynamic arrival & departure + fixed valuation
- 2 dynamic arrival & departure + dynamic valuation

# Dynamic Arrival and Departure with Fixed Valuation

- **type:**  $\theta_i = (v_i, a_i, d_i)$ 
  - $v_i$  is trader  $i$ 's (fixed) valuation of a single unit of the commodity
  - $a_i$  ( $d_i$ ) is the arrival (departure) time of trader  $i$
- an ask  $\theta_i = (v_i, a_i, d_i)$  and a bid  $\theta_j = (v_j, a_j, d_j)$  are **matchable** iff  $v_i \leq v_j$  and  $[a_i, d_i] \cap [a_j, d_j] \neq \emptyset$

## Application Example

- Stock Exchange

with Fixed Valuation

# Impossibility

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X			O				
	Maximal Matching		X						X	
with temporal constraints	Augmentation-based		X	X	X	X				
under group buying	Existence		X	X	X					
	Impossibility (cond.)		X	X	X					
	Second Price		B	X	X					
	Second Price Plus ...		S	X	X					
	Impossibility		X	X	X			X		
dynamic arrival & departure fixed valuation	Impossibility		X	X						
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based		X			O	O			O



with Fixed Valuation

# Impossibility

- There is no **deterministic** and **truthful** online double auction that can guarantee **efficiency**.

with Fixed Valuation

# Deterministic and Efficient Mechanism

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
	Maximal Matching		X						X	
with temporal constraints	Augmentation-based		X	X	X	X				
under group buying	Existence		X	X	X					
	Impossibility (cond.)		X	X	X					
	Second Price		B	X	X					
	Second Price Plus ...		S	X	X					
	Impossibility		X	X	X		X			
dynamic arrival & departure fixed valuation	Impossibility		X	X						
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based		X		O	O		O		

with Fixed Valuation

# Deterministic and Efficient Mechanism

given that:

- sellers are relatively stable
- demand is not more than supply

a *greedy* **deterministic** auction can be **truthful** and **individually rational** and guarantees **efficiency**.

with Fixed Valuation

# Deterministic and Efficient Mechanism

Environments	Mechanism		Properties							
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	Maximal Matching		X						X	
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dynamic arrival & departure fixed valuation	Impossibility		X	X						
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with Fixed Valuation

# The Reduction Framework

Environments	Mechanism		Properties							
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static	Equilibrium Matching		X				O			
	Maximal Matching		X						X	
with temporal constraints	Augmentation-based		X	X	X	X				
under group buying	Existence		X	X	X					
	Impossibility (cond.)		X	X	X					
	Second Price		B	X	X					
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	Impossibility		X	X	X		X			
dynamic arrival & departure fixed valuation	Impossibility		X	X						
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based		X		O	O		O		

# The Reduction Framework

given that:

- sellers are relatively stable
- demand is not more than the supply
- demand is known

with Fixed Valuation

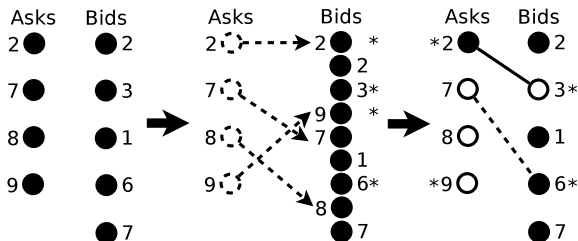
# The Reduction Framework

given that:

- sellers are relatively stable
- demand is not more than the supply
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with Fixed Valuation

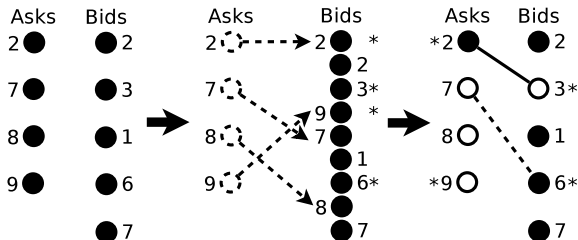
# The Reduction Framework





with Fixed Valuation

# The Reduction Framework



Properties:

- the **truthfulness** and **efficiency** of the reduced double auction **follow** that of the one-sided auction.

with Fixed Valuation

# The Reduction Framework

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
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## Dynamic Arrival, Departure and Valuation

- multiple dynamic arrival and departure
- valuation is changing over time

## Application Example

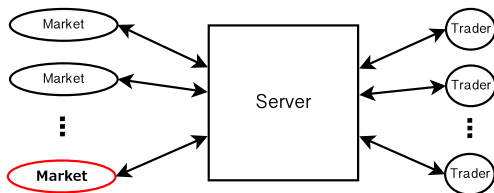
- Stock Exchange

# Dynamic Arrival, Departure and Valuation

- multiple dynamic arrival and departure
- valuation is changing over time

## Application Example

- Stock Exchange



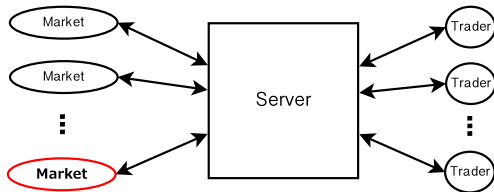
with Dynamic Valuation

# Behaviour-based Auction Design

Environments	Mechanism		Properties							
	Allocation	Payment	IC	IR	Eff.	BB	Liq.	Pro.	Com.	
static	Equilibrium Matching		X				O			
	Maximal Matching		X						X	
with temporal constraints	Augmentation-based		X	X	X	X				
under group buying	Existence		X	X	X					
	Impossibility (cond.)		X	X	X					
	Second Price		B	X	X					
	Second Price Plus ...		S	X	X					
	Impossibility		X	X	X				X	
dynamic arrival & departure fixed valuation	Impossibility		X	X						
	Greedy (cond.)		X	X	X					
	Reduction (cond.)		X	X	X					
dynamic arrival & departure dynamic valuation	Behaviour-based		X				O	O		O

# Behaviour-based Auction Design

- 1 behaviour recognition
- 2 behaviour-based classification of traders
- 3 behaviour-based mechanism design
- 4 environment adaptation



## Experimental Results

Results (based on Trading Agent Competition Market Design Tournament):

- our trading agent, *jackaroo*, achieved **1st**, **2nd** and **1st** in CAT Tournament 2009, 2010 and 2011, respectively

with Dynamic Valuation

# Experimental Results

Environments	Mechanism		Properties							
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dynamic arrival & departure fixed valuation	Impossibility		X	X						
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dynamic arrival & departure dynamic valuation	Behaviour-based		X				O	O		O



# Outline

- 1 Motivation
- 2 The Problem & Solution
- 3 Static Double Auctions
- 4 Approaching Online Double Auctions
- 5 Online Double Auctions
- 6 Conclusion
  - Summary
  - Future Work

## Summary

## Results Overview

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under group buying	Existence		X	X		X				
	Impossibility (cond.)		X	X		X				
	Second Price		B	X		X				
	Second Price Plus ...		S	X		X				
	Impossibility		X	X		X	X			
dynamic arrival & departure fixed valuation	Impossibility		X		X					
	Greedy (cond.)		X	X	X					
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dynamic arrival & departure dynamic valuation	Behaviour-based			X	O		O	O		

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# What are the potential impacts of the results?

- **Static Double Auctions**
  - more applicable allocation in stock exchange
- **Approaching Online Double Auctions**
  - augmentation-based
    - application in exchange markets, e.g. futures exchange
  - group buying related
    - a guideline for further study of the dynamic problem
- **Online Double Auctions**
  - an efficient approach to control/guide complex exchange markets

# Future Work

- 1 Group buying with dynamic advertising effect
  - how many buyers will come back?
  - will they tell their friends of the product?
- 2 How to apply the results in real applications
  - how to deal with bounded rationality of human traders?
- 3 Extending the results to other dynamic environments
  - dynamic kidney exchange
  - online advertising

# Q & A

Thank **YOU** very much  
for your attention :)