

# CS243: Introduction to Algorithmic Game Theory

Redistribution (Dengji ZHAO)

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# Recap: The General Setting of Mechanism Design

- A set of  $n$  participants/players, denoted by  $N$ .
- A mechanism needs to choose some alternative from  $A$  (allocation space), and to decide a payment for each player.
- Each player  $i \in N$  has a **private** valuation function  $v_i : A \rightarrow \mathbb{R}$ , let  $V_i$  denote all possible valuation functions for  $i$ .
- Let  $v = (v_1, \dots, v_n)$ ,  $v_{-i} = (v_1, \dots, v_{i-1}, v_{i+1}, \dots, v_n)$ .
- Let  $V = V_1 \times \dots \times V_n$ ,  $V_{-i} = V_1 \times \dots \times V_{i-1} \times V_{i+1} \times \dots \times V_n$ .

## Recap: Myerson's Optimal Auction

- Given the bids  $\mathbf{b}$  and the distribution of agents' valuations  $\mathbf{F}$ , compute **virtual bids**  $b'_i = \phi_i(b_i) = b_i - \frac{1 - F_i(b_i)}{f_i(b_i)}$ .
- Run VCG on the virtual bids  $\mathbf{b}'$  to get allocation  $\mathbf{x}'$  and payment  $\mathbf{p}'$ .
- Output  $\mathbf{x} = \mathbf{x}'$  and  $\mathbf{p}$  with  $p_i = \phi_i^{-1}(p'_i)$ .

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### Profit maximisation

Myerson's Optimal Auction maximises the seller's profit.

# Recap: Auctions

- Truthful Mechanisms
  - Second-price auction
  - Generalization: VCG auctions
  - Optimal: Myerson's mechanism
- On social networks
  - Incentive diffusion mechanism (IDM)

# Outline

## 1 Redistribution

# Alternative Objective in Auctions

- Previously we focus on **seller's revenue**.
- What if the seller is not keen on revenue (e.g., an external agent or the government)?
- We now want to return the surplus to the agents.

## Redistribution

Seeks to minimize net transfers from agents to an external body by return of VCG surplus to the agents.

# Requirements

**incentive compatibility** Each agent will truthfully report her valuation  $v_i$ .

**individual rationality** Each agent will not suffer loss when she report her true valuation.

**budget balanced** The amount of extracted wealth that cannot be redistributed among the agents is 0.



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# A First Attempt

## Question

What if we uniformly return the VCG surplus, i.e., for each agent, we return  $v_2/n$ , where  $v_2$  is the second highest bid among all agents?

# Impossibility

## Myerson-Satterthwaite Theorem

No mechanism is capable of achieving incentive compatibility, individual rationality, efficiency and budget balance at the same time.

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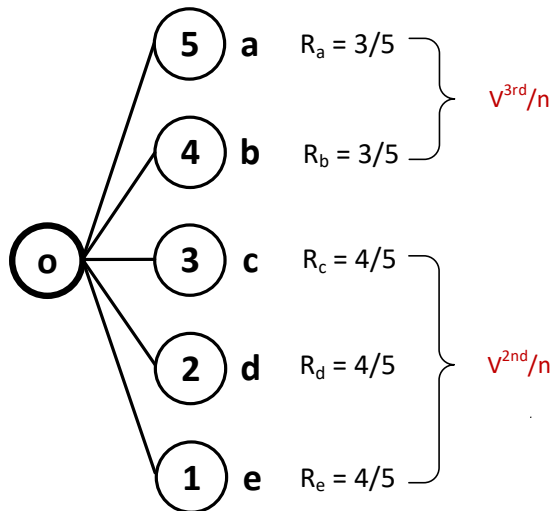
**asymptotically budget balanced** As the number of participating agents goes to infinity, the amount of extracted wealth that cannot be redistributed among the agents goes to 0.

# Cavallo's Method

- Suppose agents  $a_1, a_2, \dots, a_n$  has bids  $v'_1 \geq v'_2 \geq \dots \geq v'_n$ .
- Let  $a_1$  be the winner and pays  $v'_2$ . (VCG)
- Return the surplus  $v'_2$  back to agents as follows

$$r_i = \begin{cases} v'_3/n & \text{for } i = a_1, a_2 \\ v'_2/n & \text{for } i = a_3, \dots, a_n \end{cases}$$

# Cavallo's Method



# Cavallo's Method

The amount not redistributed is

$$r_c = v'_2 - \sum r_i = \frac{2}{n} (v'_2 - v'_3)$$

# Cavallo's Method

## Theorem

Cavallo's Method is incentive compatible, individually rational, efficient and asymptotically budget balanced.



# Advanced Reading

- *Optimal Decision Making With Minimal Waste:  
Strategyproof Redistribution of VCG Payments*  
by Ruggiero Cavallo (AAMAS 2006)

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Sponsored Search Auction (Dengji ZHAO)

SIST, ShanghaiTech University, China

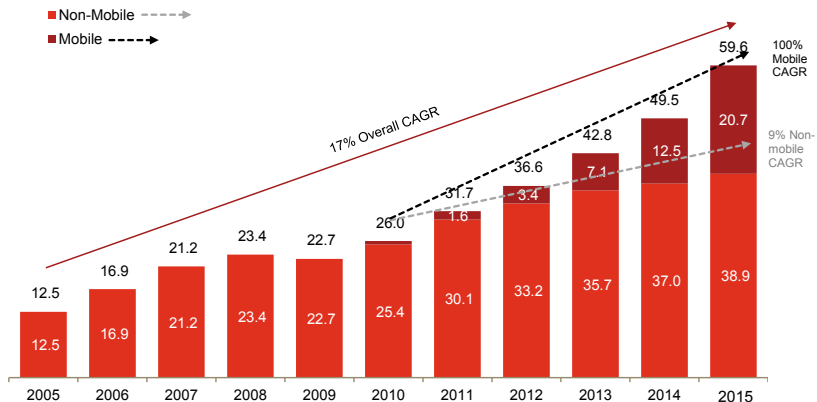
# Sponsored Search Auctions

## Sponsored Search Auction

- Used to sell ads slots by search engines such as Google, Baidu.
- Profit maximisation for the search engines?

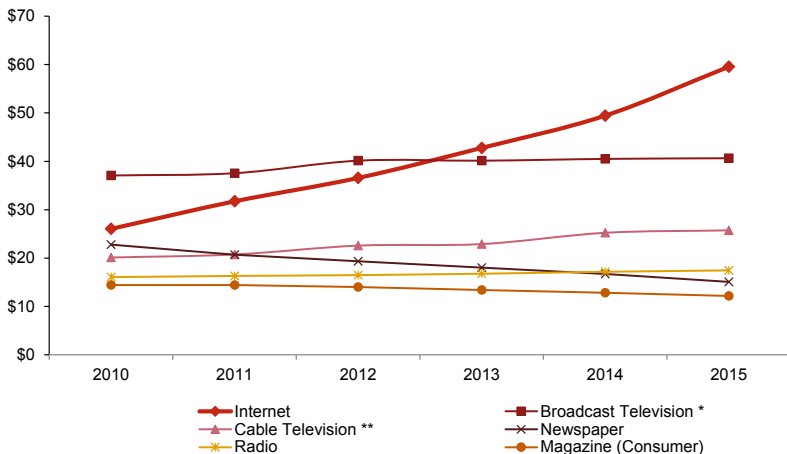
# The Sponsored Search Market

Annual Revenue 2005-2015 (\$ billions)



Source: IAB/PwC Internet Ad Revenue Report, FY 2015

# The Sponsored Search Market

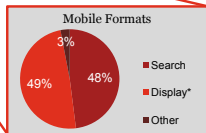
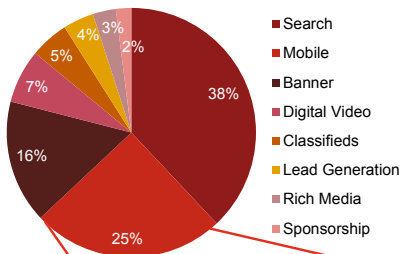


Sources: IAB/PwC Internet Ad Revenue Report, FY 2015; PwC

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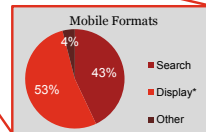
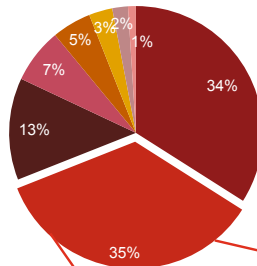
## Ad formats – full year 2014

Total - \$49.5 billion\*\*



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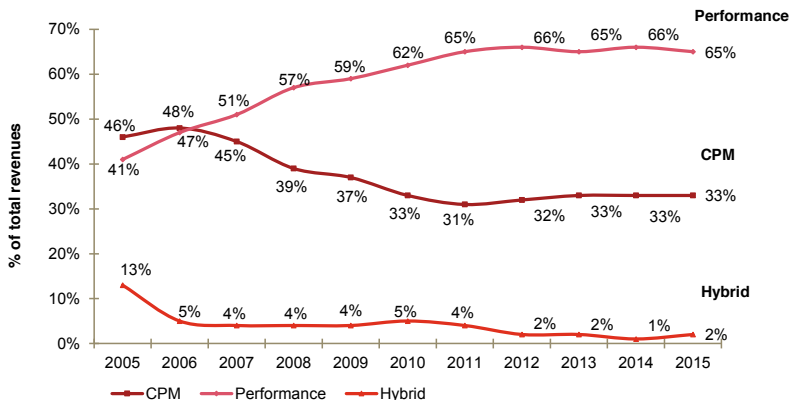
Total - \$59.6 billion\*\*



Source: IAB/PwC Internet Ad Revenue Report, FY 2015

# The Sponsored Search Market

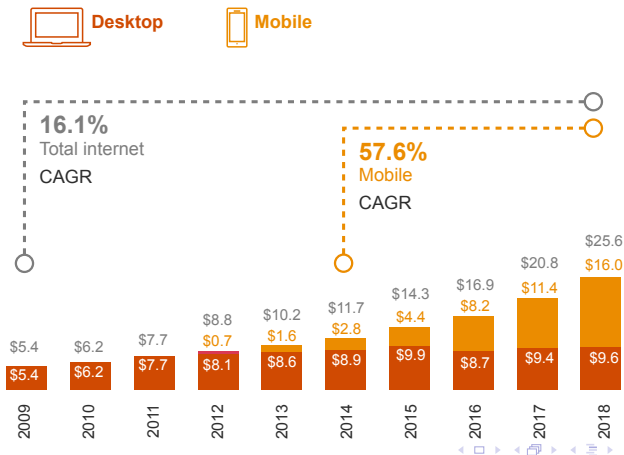
Internet ad revenues by pricing model\*



Source: IAB/PwC Internet Ad Revenue Report, FY 2015

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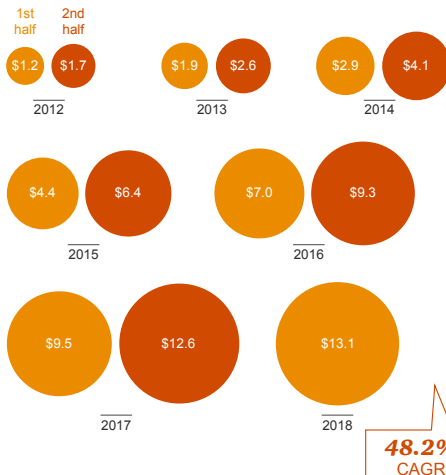
## Desktop vs. mobile internet advertising revenue (Second quarter results, \$ billions)





# The Sponsored Search Market

**Social media advertising revenue, half year results**  
(\$ billions)



Source: IAB/PwC Internet Ad Revenue Report, HY 2018

# The Basic Model

- A set of advertisers/bidders ( $n$ ), each specify a list of pairs of **keywords and bids** as well as a **total budget** (daily/weekly/monthly).
- A search engine with  $m < n$  number of **ad slots**. The search engine estimate a **click through rate**  $\alpha_{ij}$ , the probability that a user will click on the  $i$ -th slot when it is occupied by bidder  $j$ . Assume that  $\alpha_{ij} \geq \alpha_{i+1j}$  for  $i = 1, \dots, m - 1$ .
- The search engine also assigns a weight  $w_j$  to each advertiser  $j$ . The weight can be thought of as a relevance or quality metric.

# Generalized Second Price (GSP) Auctions

For each search of a keyword, GSP does the following to allocate ads:

- Rank advertisers by their **score**  $b_j w_j$ .
- The highest score gets the first slot, the second highest score gets the second slot and so on.
- A bidder pays per click the lowest bid necessary to retain his position.

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Two different variants:

- 1 Rank by bid (used by Overture): assume that  $w_j = 1$
- 2 Rank by revenue (used by Google): assume that  $w_j = \alpha_{1j}$

# Efficiency in a Static Setting

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$$\begin{aligned} \max \quad & \sum_{i=1}^k \sum_{j=1}^n \alpha_{ij} v_j x_{ij} \\ \text{s.t.} \quad & \sum_{j=1}^n x_{ij} \leq 1 \quad \forall i = 1, \dots, k \\ & \sum_{i=1}^k x_{ij} \leq 1 \quad \forall j = 1, \dots, n \\ & x_{ij} \geq 0 \quad \forall i = 1, \dots, k, \forall j = 1, \dots, n \end{aligned}$$

where  $x_{ij} = 1$  if bidder  $j$  is assigned to slot  $i$  and zero otherwise.

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- What will be the payment under VCG?

# The VCG Payments

- Consider three bidders 1, 2, 3 with  $v_1 > v_2 > v_3$  for one keyword and two slots.
- Suppose that  $\alpha_{ij} = \mu_i$  with  $\mu_1 > \mu_2$  (CTR are bidder independent).



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# Revenue Maximization

- How to maximize search engine's revenue?

# The Dynamic Setting

- What will happen if the game is repeated?

# Advanced Reading

- AGT Chapter 28. Sponsored Search Auctions